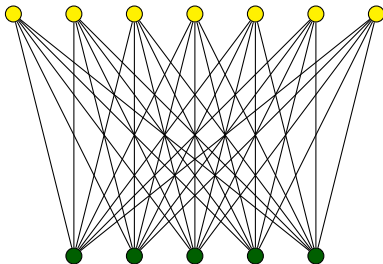


New lower bounds on crossing numbers of $K_{m,n}$

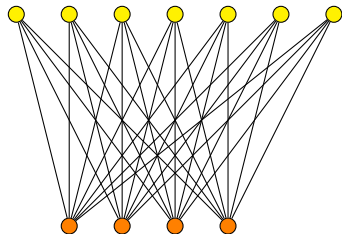
Sven Polak

CWI



Joint with Daniel Brosch (Universität Klagenfurt)
arXiv:2206.02755

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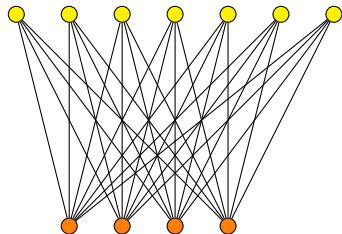


The graph $K_{7,4}$

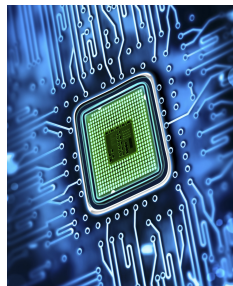
► Minimum number of crossings?

(Turán, 1940s)

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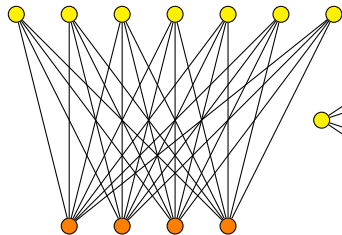


Computer chip

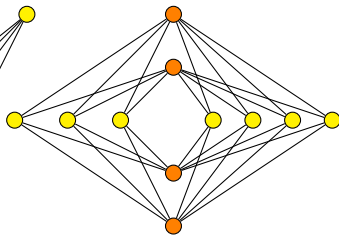
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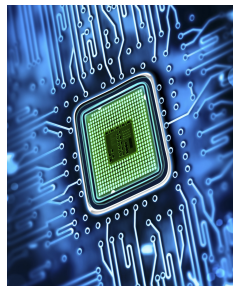
Crossing numbers of $K_{m,n}$



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Optimal drawing



Computer chip

- ▶ Minimum number of crossings?

(Turán, 1940s)

- ▶ $\text{cr}(K_{m,n}) = \lfloor \frac{(m-1)^2}{4} \rfloor \lfloor \frac{(n-1)^2}{4} \rfloor$?

(Zarankiewicz's conjecture, 1956)

Main goal

Compute lower bounds on $\text{cr}(K_{m,n})$.

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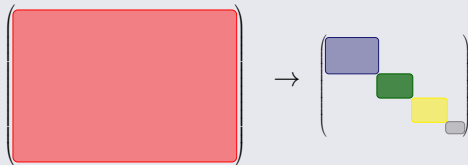
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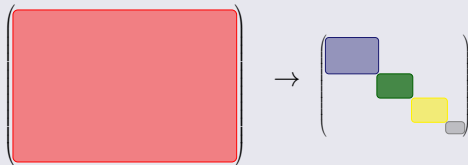
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New lower bounds on $\text{cr}(K_{n,n})$

- ▶ We also give a lower bound β_m on α_m and compute it for $m \leq 13$.

Cyclic orders

Let $\{1, \dots, m\}$ and $\{b_1, \dots, b_n\}$ be the sides of the bipartition of $K_{m,n}$. Given a drawing of $K_{m,n}$, define for each vertex b_i :

$\gamma(b_i) :=$ the cyclic order in which edges from b_i go to $\{1, \dots, m\}$

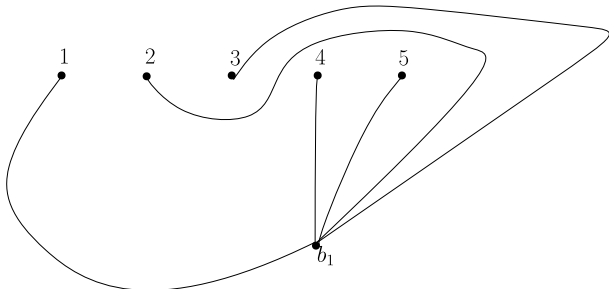


Figure: $\gamma(b_1) = (14523)$.

Lemma (Kleitman; 1970, Woodall; 1993)

#crossings of edges incident with b_i and $b_j \geq$
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Example

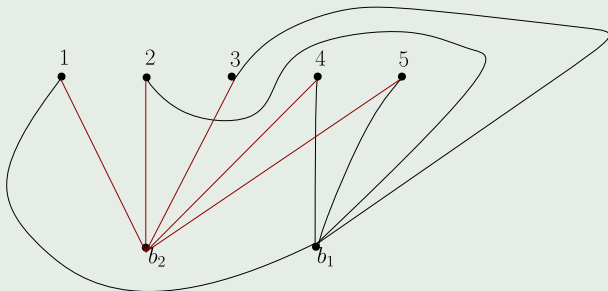


Figure: $\gamma(b_2) = (12345)$ and $\gamma(b_1)^{-1} = (14523)^{-1} = (13254)$.

Min. #swaps of adjacent elements of $\gamma(b_2)$ to change $\gamma(b_2)$ into $\gamma(b_1)^{-1}$ is 2.

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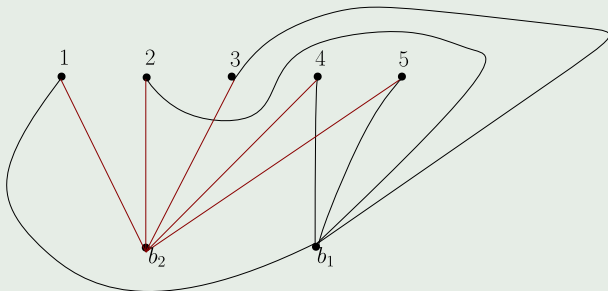


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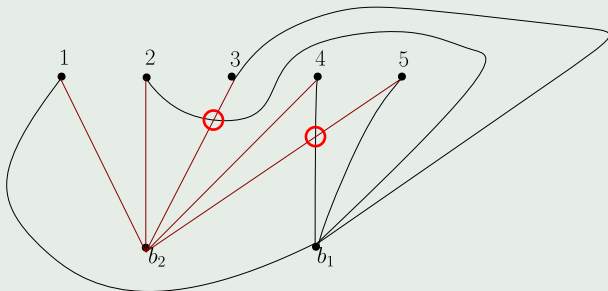


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Lower bound on crossing number

Let Z_m be the set of m -cycles, and let $Q \in \mathbb{R}^{Z_m \times Z_m}$ with

$Q_{\sigma, \tau} := \min. \# \text{swaps of adjacent elements of } \sigma \text{ to change } \sigma \text{ into } \tau^{-1}.$

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Define $q_m := \min \left\{ x^T Q x \mid x \in \mathbb{R}_{\geq 0}^{Z_m}, \mathbf{1}^T x = 1 \right\}$. Then

$$\text{cr}(K_{m,n}) \geq \frac{1}{2} n^2 q_m - \frac{1}{2} n \lfloor \frac{1}{4} (m-1)^2 \rfloor \text{ for all } m, n.$$

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Semidefinite relaxation α_m of q_m

$q_m \geq \alpha_m := \min \{ \langle X, Q \rangle \mid X \in \mathbb{R}_{\geq 0}^{Z_m \times Z_m}, \langle X, J \rangle = 1, X \text{ positive semidefinite} \}.$

Exploiting symmetry for computing α_m

Let $G_m := S_m \times \{-1, +1\}$. Then G_m acts on Z_m via

$$(\pi, \varepsilon) \cdot \sigma = \pi \sigma^\varepsilon \pi^{-1}, \quad \text{for } (\pi, \varepsilon) \in G_m \text{ and } \sigma \in Z_m.$$

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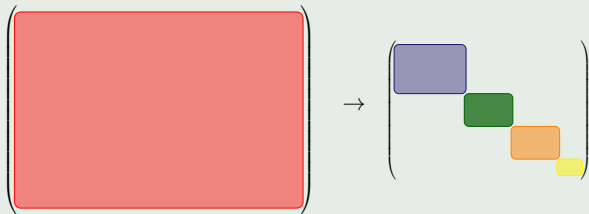
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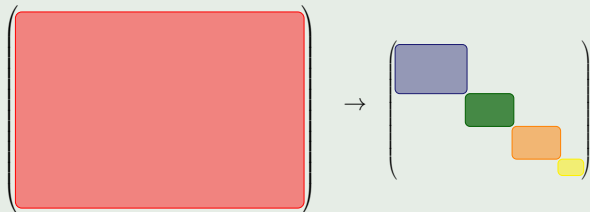
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m	$(m-1)!$	$\sum m_i^2$	max block size
7	720	78	3
8	5040	380	7
9	40320	2438	12
10	362880	18744	38

Group invariance and Artin-Wedderburn

Let G be a finite group acting on a finite set Z . Decompose

$$\mathbb{C}^Z = \bigoplus_{i=1}^k \bigoplus_{j=1}^{m_i} V_{i,j},$$

for irreducible G -modules $V_{i,j}$ with $V_{i,j} \cong V_{i',j'}$ iff $i = i'$.

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$$\begin{aligned} \Phi : \left(\mathbb{C}^{Z \times Z}\right)^G &\rightarrow \bigoplus_{i=1}^k \mathbb{C}^{m_i \times m_i}, \\ A &\mapsto \bigoplus_{i=1}^k U_i^* A U_i, \quad \text{where } U_i := (u_{i,j} \mid j \in [m_i]). \end{aligned}$$

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Key fact

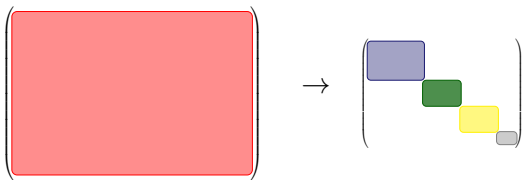
For all $A \in \left(\mathbb{C}^{Z \times Z}\right)^G$ we have $A \succeq 0 \iff \Phi(A) \succeq 0$.

Surjective S_m -homomorphism

$$f : M^{(1^m)} \rightarrow \mathbb{C}^{Z_m},$$

$$\begin{pmatrix} \overline{i_1} \\ \overline{i_2} \\ \vdots \\ \overline{i_m} \end{pmatrix} \mapsto \text{indicator vector in } \mathbb{C}^{Z_m} \text{ of } (i_1 i_2 \dots i_m).$$

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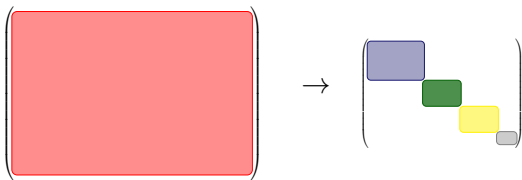
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► See thesis Daniel Brosch: faster decomposition of modules M^μ / F .

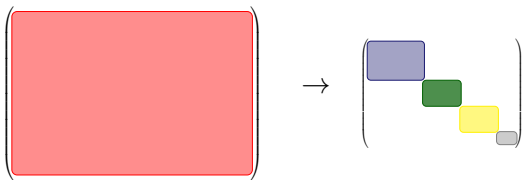


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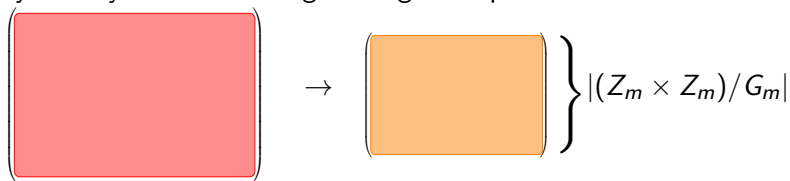
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3. Reduce further with **inversion** ($\{\pm 1\}$)-action, splits each block in 2.



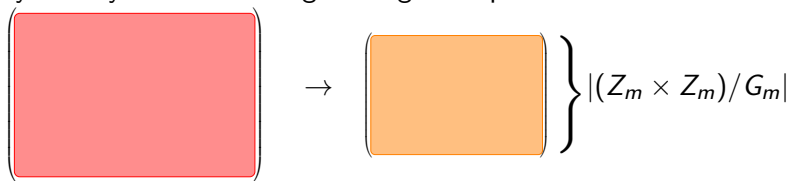
Results

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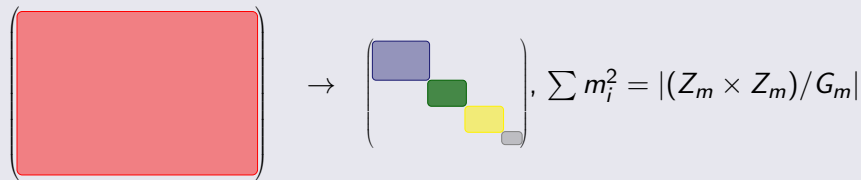


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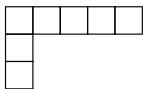


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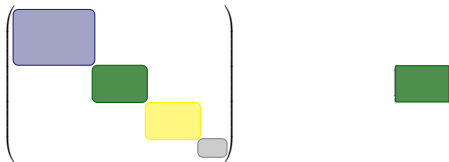
The relaxation β_m



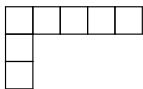
$$\lambda = (m - 2, 1, 1)$$

Observation

Requiring only one specific block of size $\lfloor \frac{m-1}{2} \rfloor$ to be PSD gives a good lower bound on α_m .



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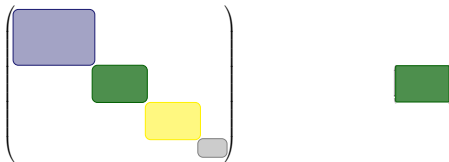


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- ▶ Dual of β_m : **few variables**, **small matrix block**, **many linear constraints**.
- ▶ We solve the program without constraints, and iteratively add the most violated constraint until all constraints are satisfied.



Theorem (Brosch, P.; 2022+)

$$\text{cr}(K_{10,n}) \geq 4.8345n^2 - 10n,$$

$$\text{cr}(K_{11,n}) \geq 5.9088n^2 - 12.2222n,$$

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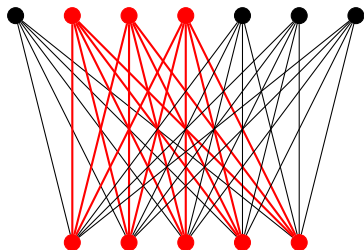
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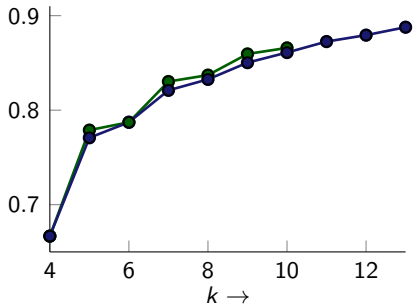
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Asymptotic results

Lemma (de Klerk, Maharry, Pasechnik, Richter, Salazar; 2006)

$$\lim_{n \rightarrow \infty} \frac{\text{cr}(K_{n,n})}{Z(n,n)} \geq \frac{8\alpha_k}{k(k-1)}.$$

k	$\frac{8\alpha_k}{k(k-1)}$	$\frac{8\beta_k}{k(k-1)}$
7	0.8303	0.8210
8	0.8371	0.8326
9	0.8595	0.8503
10	0.8659	0.8610
11		0.8726
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13		0.8878



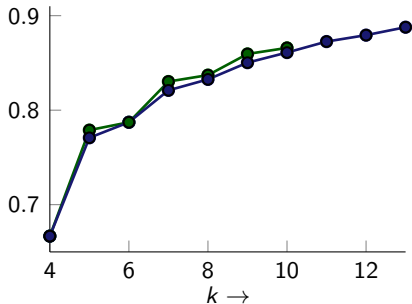
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Lemma (Brosch, P.; 2022+)

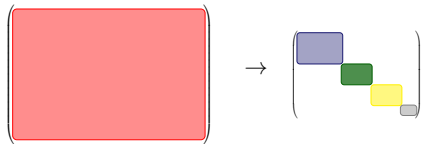
$$\lim_{n \rightarrow \infty} \frac{\text{cr}(K_{m,n})}{Z(m,n)} \geq 0.8878 \frac{m}{m-1} \text{ for each } m \geq 13.$$

Summary

- ▶ Semidefinite programming gives good lower bounds on $\text{cr}(K_{m,n})$.

n	best lower bound previously known	new lower bound	$Z(n, n)$
10	384	388	400
11	581	589	625
12	846	865	900
13	1192	1229	1296

New lower bounds on $\text{cr}(K_{n,n})$

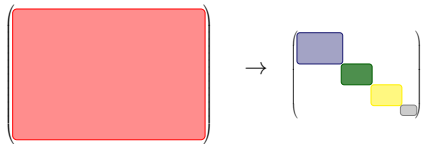


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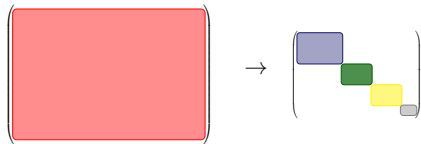
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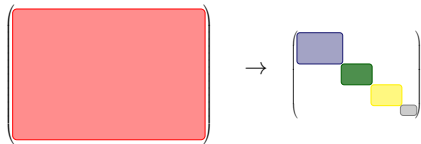


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- ▶ Possible to give optimum solution to β_m analytically?

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The relaxation β_m : structure in the solutions?

The dual of β_m

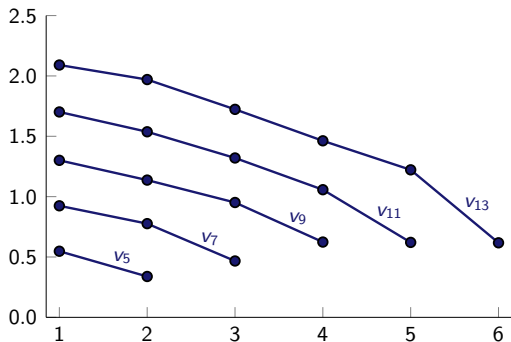
$$\beta_m = \max \left\{ t : Y \in \mathbb{R}^{\lfloor \frac{m-1}{2} \rfloor \times \lfloor \frac{m-1}{2} \rfloor}, Y \succeq 0, \right. \\ \left. \forall \omega \in (Z_m \times Z_m)/G_m : \langle Y, A_\omega \rangle + |\omega|t \leq |\omega|q_\omega \right\},$$

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where $q_\omega := Q_{\sigma,\tau}$ for any $(\sigma,\tau) \in \omega$, and $A_\omega := U_\lambda^T \left(\sum_{(\sigma,\tau) \in \omega} E_{\sigma,\tau} \right) U_\lambda$.



Vectors v_m such that optimal $Y = \frac{1}{(m-1)!} v_m v_m^T$.

Lower bound on crossing number

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$$Q_{\sigma, \tau} := \min. \# \text{swaps of adjacent elements to change } \sigma \text{ into } \tau^{-1}.$$

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$$\begin{aligned} n^2 q_m &\leq n^2 x^T Q x = c^T Q c = \sum_{\sigma, \tau \in Z_m} c_\sigma c_\tau Q_{\sigma, \tau} = \sum_{i, j=1}^n Q_{\gamma(b_i), \gamma(b_j)} \\ &= 2 \sum_{i < j} Q_{\gamma(b_i), \gamma(b_j)} + \sum_{i=1}^n Q_{\gamma(b_i), \gamma(b_i)} \leq 2 \text{cr}(K_{m,n}) + n \lfloor \frac{1}{4}(m-1)^2 \rfloor. \quad \square \end{aligned}$$