Apache Graph Processing on Hadoop

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Graphs are simple



A computer network



A social network



A semantic network



President of the United States

Hawaii

A map



Predicting break ups



Graph approach

Aggregation approach

Graphs are nasty.

Each vertex depends on its neighbours, recursively.

Recursive problems are nicely solved iteratively.



PageRank in MapReduce

- Record: < v_i, pr, [v_j, ..., v_k] >
- Mapper: emits < v_j, pr / #neighbours >
- Reducer: sums the partial values

MapReduce dataflow



Drawbacks

- Each job is executed N times
- Job bootstrap
- Mappers send PR values and structure
- Extensive IO at input, shuffle & sort, output



Timeline

- Inspired by Google Pregel (2010)
- Donated to ASF by Yahoo! in 2011
- Top-level project in 2012
- 1.0 release in January 2013
- 1.1 release in November 2014

Plays well with Hadoop















Code

def compute(vertex, messages):

minValue = Inf # float('Inf')

for m in messages:

minValue = min(minValue, m)

if minValue < vertex.getValue():</pre>

vertex.setValue(minValue)

for edge in vertex.getEdges():

```
message = minValue + edge.getValue()
```

```
sendMessage(edge.getTargetId(), message)
```

```
vertex.voteToHalt()
```





Superstep 0:

- 1. All vertices initialise their value at random.
- 2. User vertices send their value to their endpoints.
- 3. All vertices vote to halt.

The item vector updates its latent vector based on the latent vectors coming from the users, and spreads it around.

Superstep 1:

- 1. Item vertices are woken up.
- Item vertices update their values according to SGD based on the messages, edge values and the current latent vector
- 3. Item vertices send their new value to their endpoints.
- 4. Item vertices vote to halt.

It's now the turn of the user vectors to update their latent vectors based on the items' latent vectors computed in the previous super step

Superstep 2:

- 1. User vertices are woken up.
- User vertices update their values according to SGD based on the messages, edge values and the current latent vector
- 3. User vertices send their new value to their endpoints.
- 4. User vertices vote to halt.

Inactive item vertex with latent vector



(10

(16)

- 2. Edge values are initialised to these labels.
- 3. At this first superstep, all frequencies are 1 so vertices break ties randomly.
- 4. Vertices update their value and send it to the endpoints of their edges ..

The central vertex acquires at random one of the lanes from the neighbours



BSP & Giraph PU I PU 2 PU 3 PU 4 PU 5 Iteration i Iteration i+1

Advantages

- No locks: message-based communication
- No semaphores: global synchronization
- Iteration isolation: massively parallelizable

Designed for iterations

- Stateful (in-memory)
- Only intermediate values (messages) sent
- Hits the disk at input, output, checkpoint
- Can go out-of-core

Giraph job lifetime



Architecture



Composable API







No SPoFs



Giraph scales



ref: https://www.facebook.com/notes/facebook-engineering/scaling-apache-giraph-to-a-trillionedges/10151617006153920

Giraph is fast

- 100x over MR (Pr)
- jobs run within minutes
- given you have resources
 ;-)



Serialised objects



Primitive types

- Autoboxing is expensive
- Objects overhead (JVM)
- Use primitive types on your own
- Use primitive types-based libs (e.g. fastutils)

Sharded aggregators





Okapi

- Apache Mahout for graphs
- Graph-based recommenders: ALS, SGD, SVD++, etc.
- Graph analytics: Graph partitioning, Community Detection, K-Core, etc.





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