

Roberto De Virgilio

Antonio Maccioni

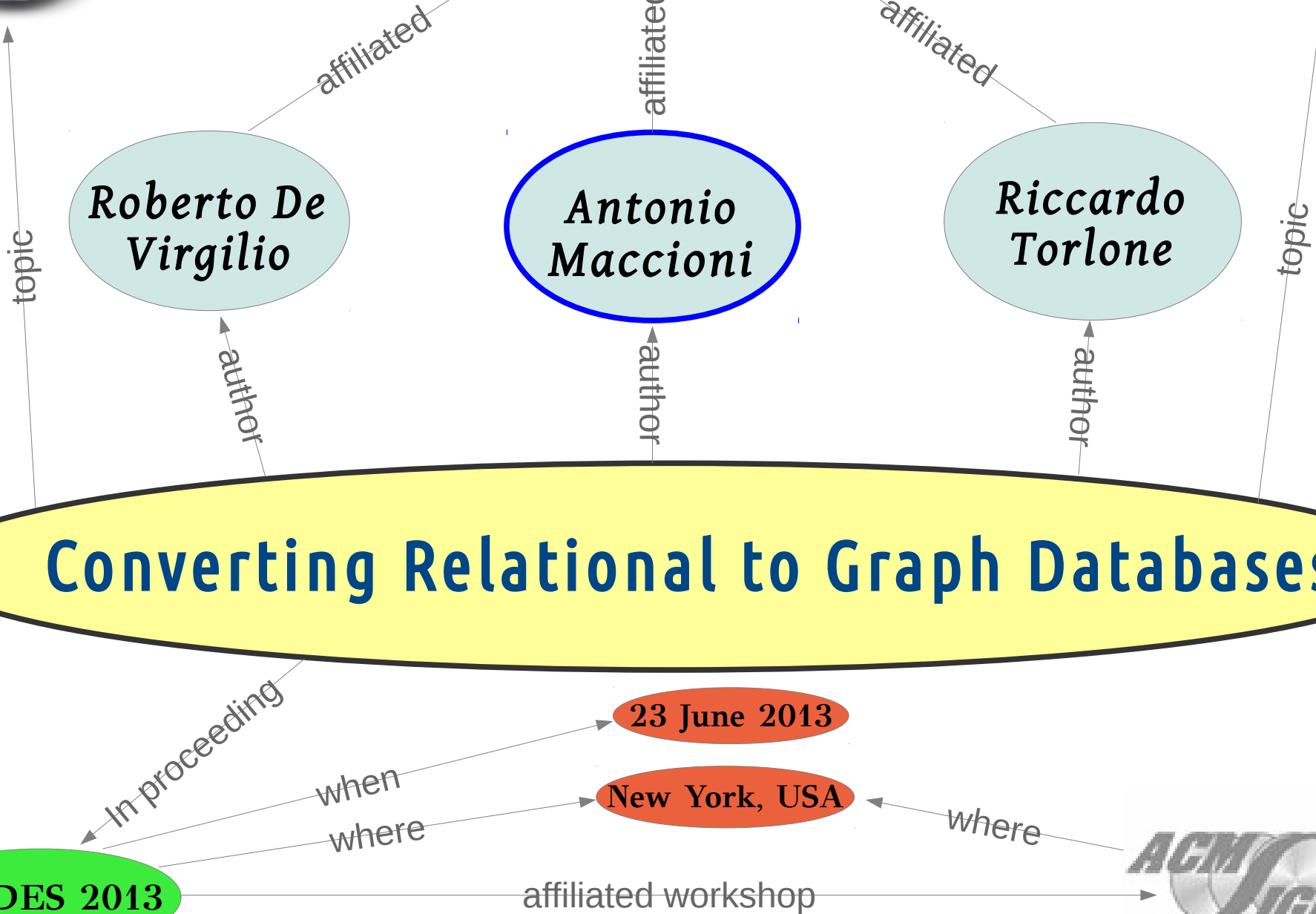
Riccardo Torlone

Converting Relational to Graph Databases

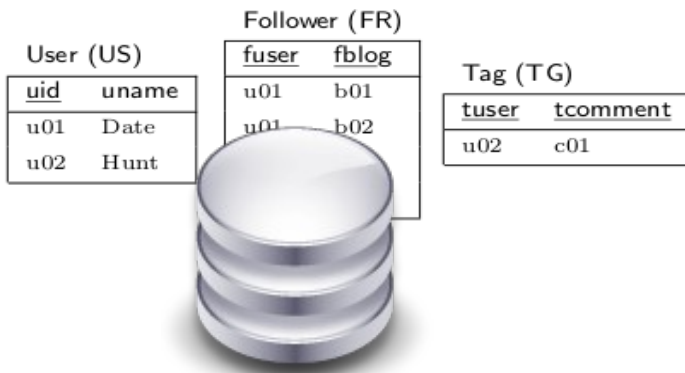
GRADES 2013

23 June 2013

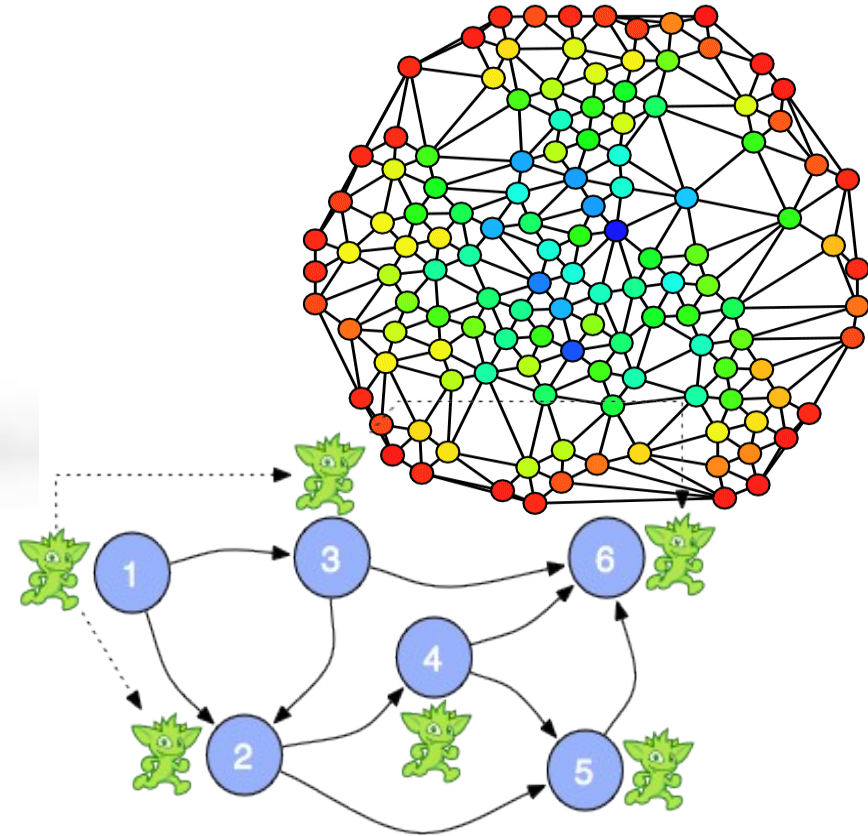
New York, USA



Relational Database Migration

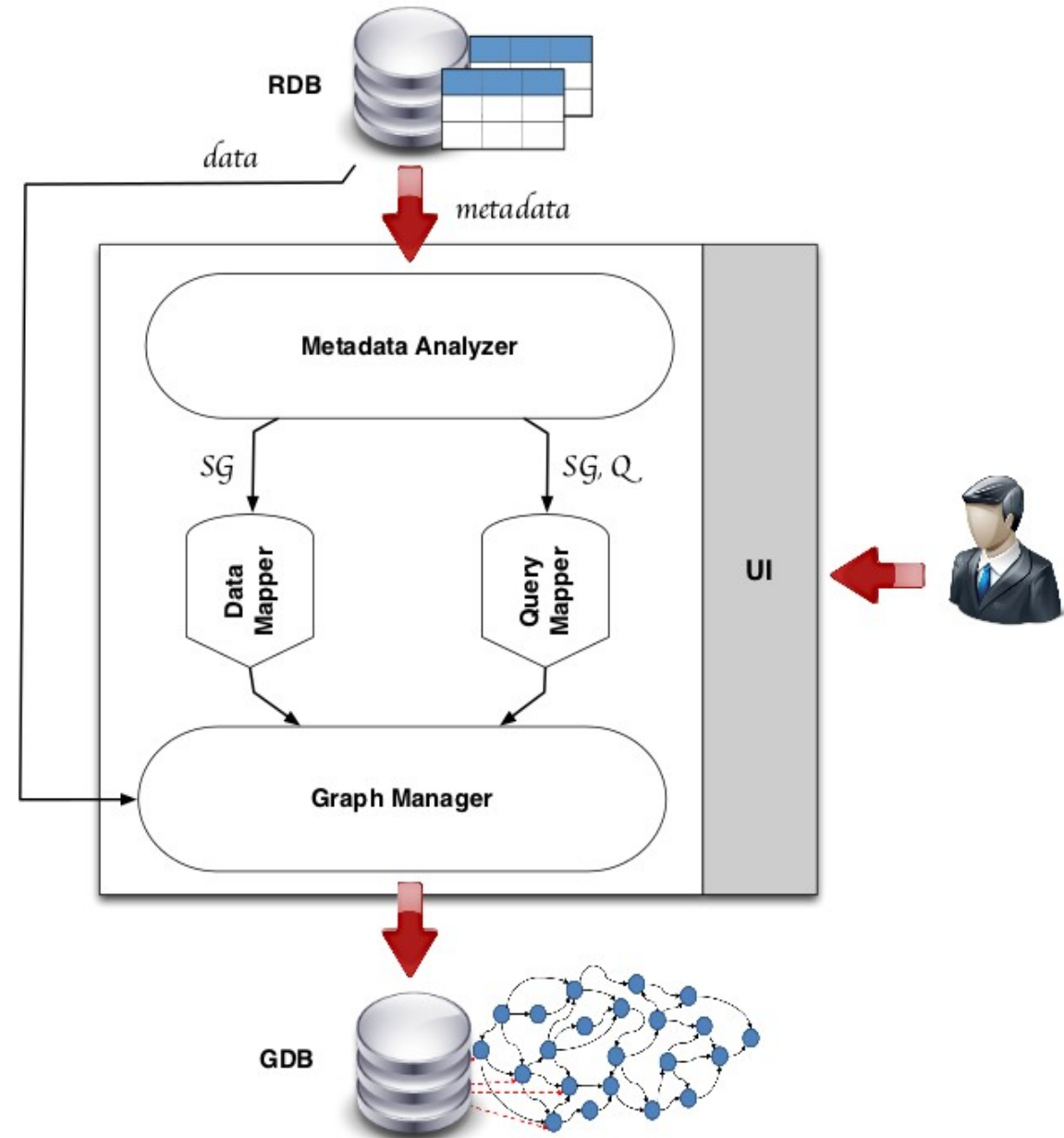


SQL
select *
from T
where T.A1 = v1



R2G: Features

- **Data** migration
- **Query** translation
- **Automatic** non-naïve approach
- Try to **minimize** the memory accesses



Graph Modeling of Relational DB

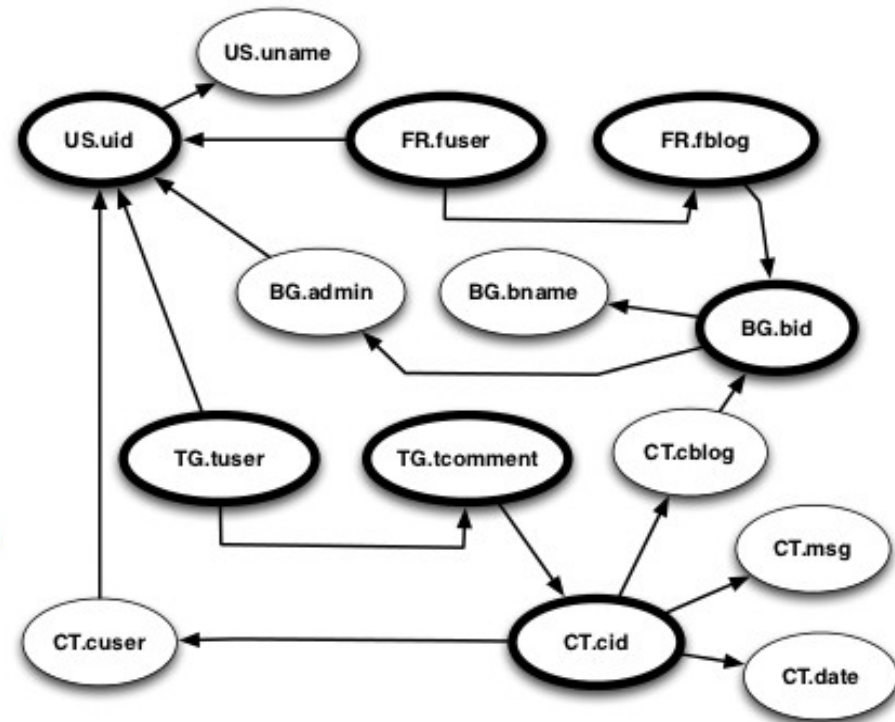
<u>uid</u>	uname
u01	Date
u02	Hunt

<u>fuser</u>	<u>fblog</u>
u01	b01
u01	b02
u01	b03
u02	b01

<u>tuser</u>	<u>tcomment</u>
u02	c01

<u>bid</u>	bname	admin
b01	Information Systems	u02
b02	Database	u01
b03	Computer Science	u02

<u>cid</u>	<u>cblog</u>	<u>cuser</u>	msg	date
c01	b01	u01	Exactly what I was looking for!	25/02/2013



- Full Schema Paths:

FR.fuser → US.uid → US.uname

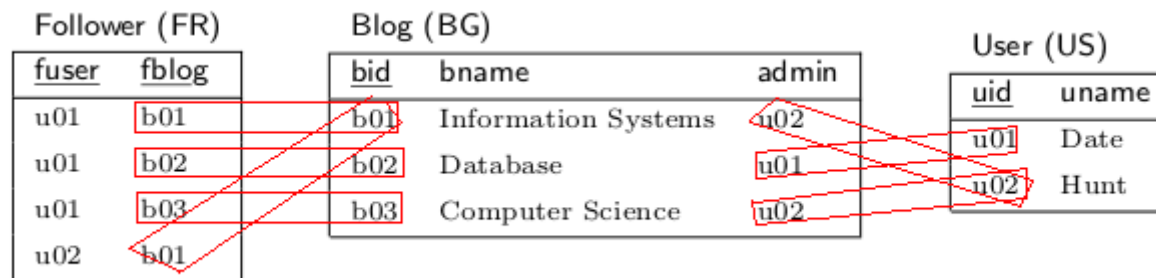
FR.fuser → FR.fblog → BG.bid → BG.bname

FR.fuser → FR.fblog → BG.bid → BG.admin → US.uid → US.uname

...

Basic Concepts

- **Joinable** tuples $t_1 \in R_1$ and $t_2 \in R_2$:
 - there is a foreign key constraint between $R_1.A$ and $R_2.B$ and $t_1[A] = t_2[B]$.
- **Unifiability** of data values $t_1[A]$ and $t_2[B]$:
 - (i) $t_1=t_2$ and both A and B do not belong to a multi-attribute key;
 - (ii) t_1 and t_2 are joinable and A belongs to a multi-attribute key;
 - (iii) t_1 and t_2 are joinable, A and B do not belong to a multi-attribute key and there is no other tuple t_3 that is joinable with t_2 .



Data Migration (1)

- Identify unifiable data exploiting **schema** and **constraints**



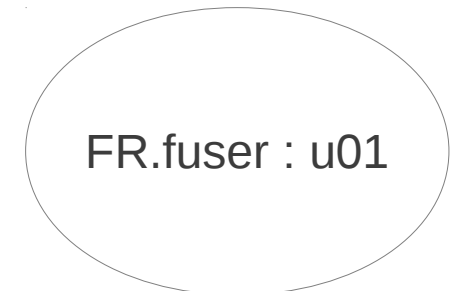
Follower (FR)

<u>fuser</u>	<u>fblog</u>
u01	b01
u01	b02
u01	b03
u02	b01

User (US)

<u>uid</u>	<u>uname</u>
u01	Date
u02	Hunt

n1



Data Migration (2)

- Identify unifiable data exploiting **schema** and **constraints**



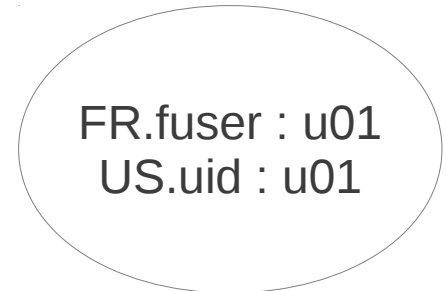
Follower (FR)

<u>fuser</u>	<u>fblog</u>
u01	b01
u01	b02
u01	b03
u02	b01

User (US)

<u>uid</u>	uname
u01	Date
u02	Hunt

n1



Data Migration (3)

- Identify unifiable data exploiting **schema** and **constraints**



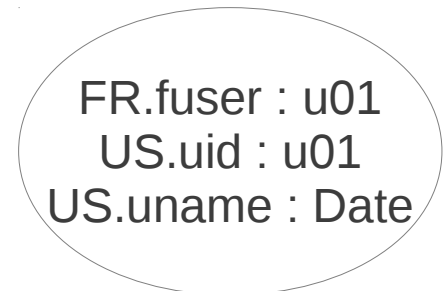
Follower (FR)

<u>fuser</u>	<u>fblog</u>
u01	b01
u01	b02
u01	b03
u02	b01

User (US)

<u>uid</u>	<u>uname</u>
u01	Date
u02	Hunt

n1



Data Migration (4)

- Identify unifiable data exploiting **schema** and **constraints**

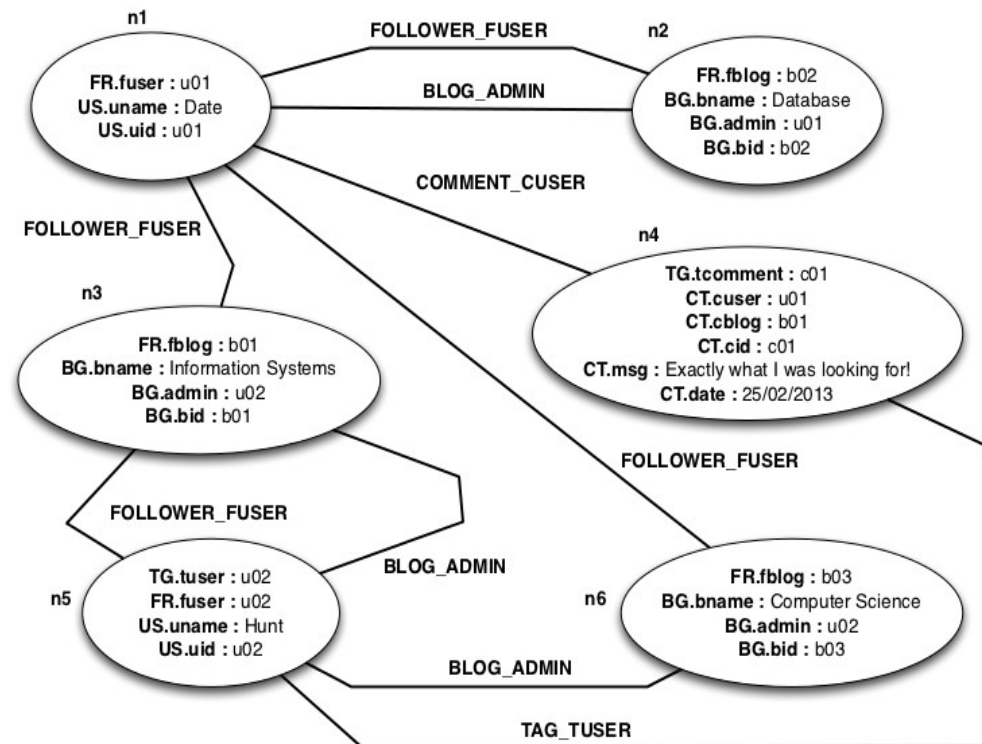
uid	uname
u01	Date
u02	Hunt

fuser	fblog
u01	b01
u01	b02
u01	b03
u02	b01

tuser	tcomment
u02	c01

bid	bname	admin
b01	Information Systems	u02
b02	Database	u01
b03	Computer Science	u02

cid	cblog	cuser	msg	date
c01	b01	u01	Exactly what I was looking for!	25/02/2013



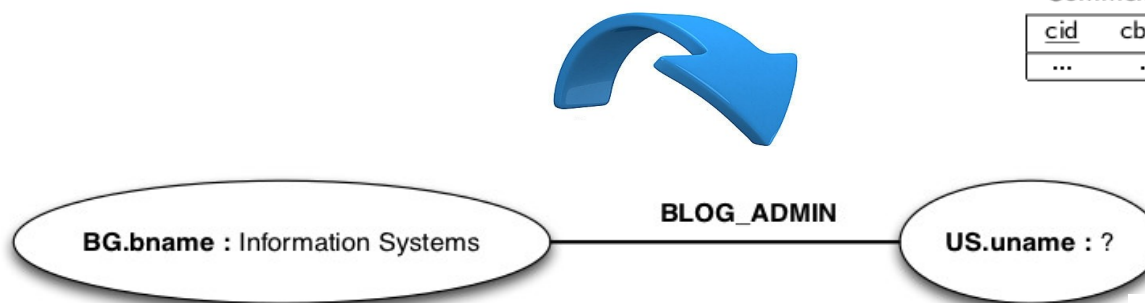
Query Translation

```
select  US.uname
from    User US, Blog BG
where   (BG.admin = US.uid) and
        (BG.bname = 'Inf. Systems')
```

User (US)		Follower (FR)		Tag (TG)	
uid	uname	fuser	fblog	tuser	tcomment
...

Blog (BG)		
bid	bname	admin
...

Comment (CT)				
cid	cblog	cuser	msg	date
...



Gremlin

```
g.V.filter{it.BG.bname=='Inf. Systems'}.
  .outE.filter{it.label=='BLOG_ADMIN'}.
  .inV.US.uname
```

XQuery

```
for  $x in /[BG.bname='Inf. Systems'],
     $y in $x/BLOG_ADMIN/*
return $y/US.uname
```

Experimental Results

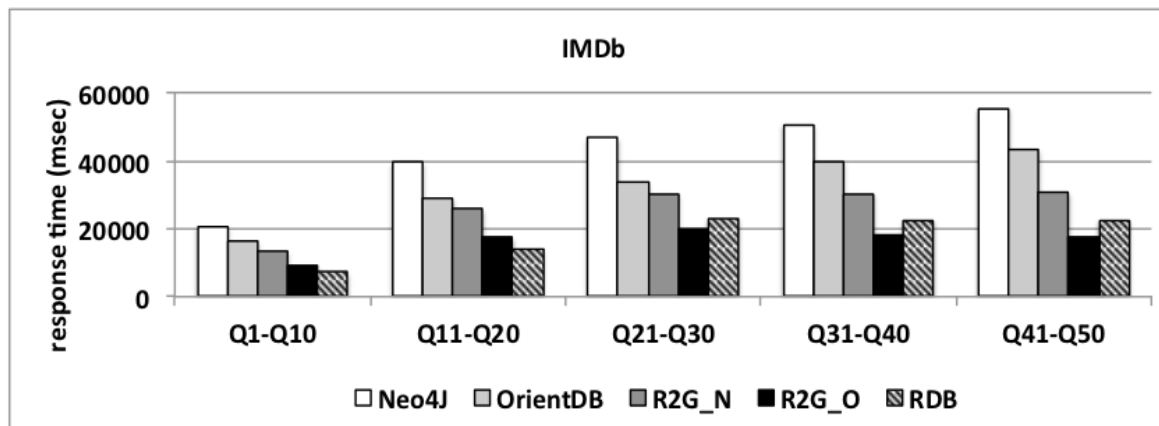
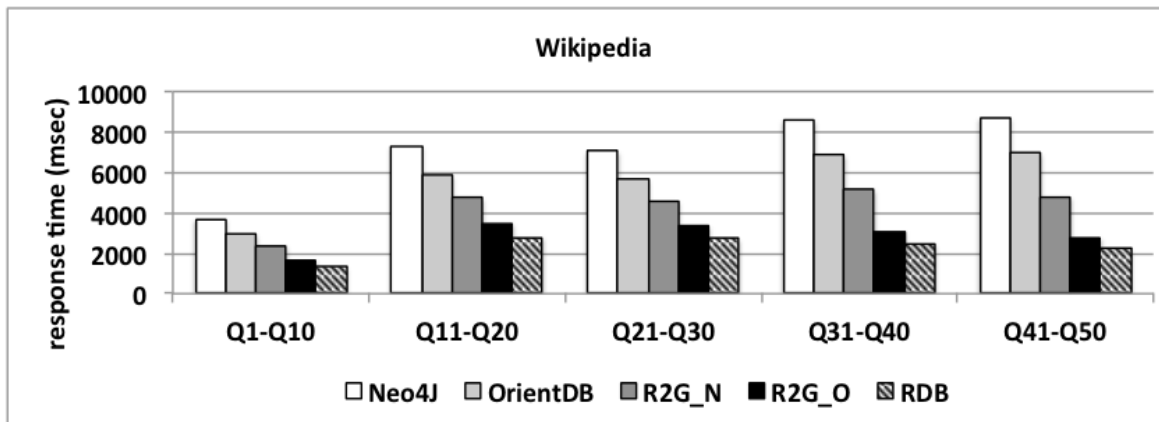
Dataset	Neo4J	OrientDB	R2G_N	R2G_O
MONDIAL	7.4 sec	5.3 sec	13.9 sec	9.3 sec
WIKIPEDIA	70.7 sec	66.5 sec	161.5 sec	148.7 sec
IMDB	8.1 min	10.2 min	16.2 min	22.1 min



The
MONDIAL
Database



WIKIPEDIA



Gremlin
 $G = (V, E)$

Conclusion

- **Automatic** data mapping
- Conjunctive **query translation** into a path traversal query
- **Independent** from a specific GDBMS
- **Efficient** exploitation of Graph Database Features

Future Work

- Consider **frequent queries** to migrate data
- Consider wider range of queries than CQ
- Improve **compactness** of the graph database

Thanks For The Attention



... **demo** presentation during the following **interactive** session!